

## CS256-Assignment#3

### 3 Problems. 50 points.

Finish reading Chapter 0 of the textbook (excluding Section 0.4 and the parts related to grouped statements).

1. [30 points] Prove that every formula of the form

$$\nabla_1 \nabla_2 \dots \nabla_n p,$$

where each  $\nabla_i$  is either  $\square$  or  $\diamond$ , is congruent to one among

$$\square p, \quad \diamond p, \quad \square \diamond p, \quad \text{and} \quad \diamond \square p.$$

For example,  $\diamond \square \square \diamond \diamond \diamond p \approx \square \diamond p$ .

Describe a procedure for computing the correct congruence.

Note: Prove the congruences, such as  $\square \square p \approx \square p$ , used in your proofs.

2. [10 points] Prove that:  $\widehat{\mathcal{W}}, \wedge, \neg, \forall$  is a complete set of operators for future temporal formulas, i.e., every future temporal formula is congruent to a formula built from state formulas using only  $\widehat{\mathcal{W}}, \wedge, \neg,$  and  $\forall$ .

Clarification:  $\forall$  is only included because temporal logic formulas include the quantifiers  $\forall$  and  $\exists$ . Don't use  $\forall$  when showing how to rewrite a temporal operator (e.g., by just opening up its definition). [Give the temporal congruences without proof.]

3. [10 points]

(a) Write a quantifier-free formula stating that  $p$  holds precisely at all even positions, i.e.,  $p$  is true at positions  $0, 2, 4, \dots$  and false at positions  $1, 3, 5, \dots$ .

(b) Using flexible quantification over boolean variables, write a formula stating that  $p$  holds at all even positions. The formula should not restrict the value of  $p$  at odd positions.

(c) Using flexible quantification over integer variables, write a formula stating that  $p$  holds at positions  $0, 1, 4, 9, 16, \dots$ , i.e., positions  $j = k^2$  for  $k = 0, 1, \dots$ . Nothing is said about the value of  $p$  at other positions. You can also use rigid quantification if necessary.